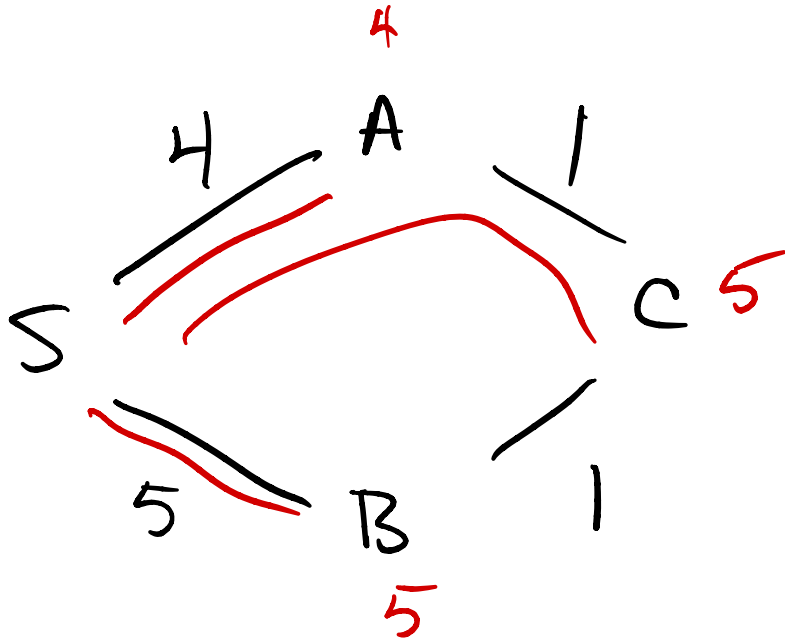


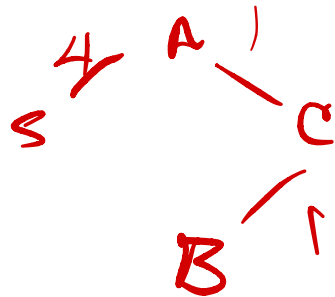
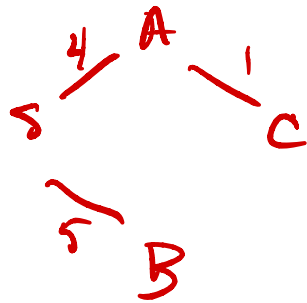
CS 252

M, 13 May 2024

8c



DA.



7.

① Given $e = (u, v)$ with weight w

① BFS in T from u to
find shortest path

$O(|V|)$

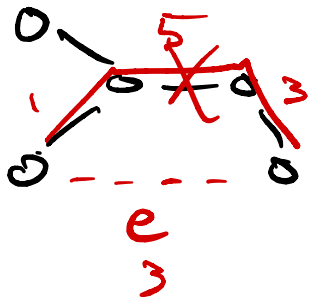
$$u = u_1, u_2, \dots, u_x = v$$

② Find max edge weight in that
path (u_i, u_{i+1}) has weight w'

$O(|V|)$

③ if $w' > w$, delete (u_i, u_{i+1})
+ add (u, v)
otherwise, leave e out

$O(1)$

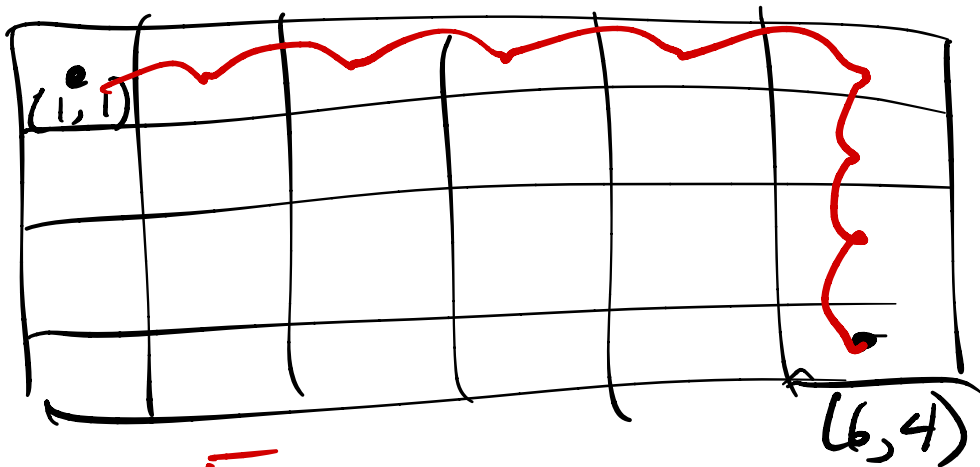


$M(6,4)$

6

Walk $(1,1)$ to
 $(6,4)$ only
R + Down

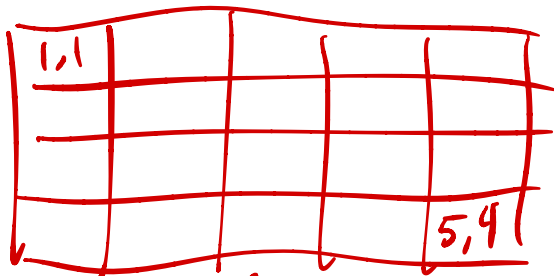
4



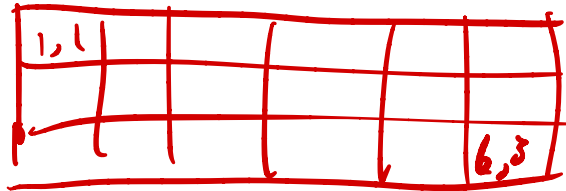
How many
such
routes

5

4



$M(5,4)$



$M(6,3)$

	size	value
0		
1	2	7
2	5	6
3	1	1
4	1	2
5	3	4
6	1	9
7	8	10
8	6	6

	d												
	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0												
2	0												
3	0												
4	0												
5	0												
6	0												
7	0												
8	0												

$$M(k, d) = \max \begin{cases} \text{value}[k] + M(k-1, d - \text{size}[k]) & \text{if } k > 0, d \geq \text{size}[k] \\ M(k-1, d) & \text{if } k > 0, d \geq 0 \\ 0 & \end{cases}$$

	size	value
0		
1	2	7
2	5	6
3	1	1
4	1	2
5	3	4
6	1	9
7	8	10
8	6	6

	d												
	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	7	7	7	7	7	7	7	7	7	7	7
2	0	0											
3	0	1											
4	0	2											
5	0	2											
6	0	9											
7	0	9											
8	0	9											

$$M(k, d) = \max \begin{cases} \text{value}[k] + M(k-1, d - \text{size}[k]) & \text{if } k > 0, d \geq \text{size}[k] \\ M(k-1, d) & \text{if } k > 0, d \geq 0 \\ 0 & \end{cases}$$

	size	value
0	X	X
1	2	7
2	5	6
3	1	1
4	1	2
5	3	4
6	1	9
7	8	10
8	6	6

	d												
	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	7	7	7	7	7	7	7	7	7	7	7
2	0	0	7	7	7	7	7	13	13	13	13	13	13
3	0	1											
4	0	2											
5	0	2											
6	0	9											
7	0	9											
8	0	9											9

goal answer

$$M(k, d) = \max \begin{cases} \text{value}[k] + M(k-1, d - \text{size}[k]) & \text{if } k > 0, d \geq \text{size}[k] \\ M(k-1, d) & \text{if } k > 0, d \geq 0 \\ 0 & \end{cases}$$

Algorithm

$M(n, C)$

$$O(n) \left\{ \begin{array}{l} \text{for } k = 0 \text{ to } n \\ M[k, 0] = 0 \end{array} \right.$$

$$O(C) \left\{ \begin{array}{l} \text{for } d = 0 \text{ to } C \\ M[0, d] = 0 \end{array} \right.$$

$$O(nC) \left\{ \begin{array}{l} \text{for } k = 1 \text{ to } n \\ \text{for } d = 1 \text{ to } C \\ m_1 = \text{value}[k] + M[k-1, d - \text{size}[k]] \text{ if } d \geq \text{size}[k] \\ \text{else } 0 \\ m_2 = M[k-1, d] \\ M[k, d] = \max(m_1, m_2) \end{array} \right.$$

		d												
		0	1	2	3	4	5	6	7	8	9	10	11	12
k	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	1	0	0	7	7	7	7	7	7	7	7	7	7	7
	2	0	0	7	7	7	7	7	13	13	13	13	13	13
	3	0	1	7	8	8	8	8	13	14	14	14	14	14
	4	0	2	7	9	10	10	10	13	15	16	16	16	16
	5	0	2	7	9	10	11	13	14	15				
	6	0	9											
	7	0	9											
	8	0	9											

$$M(k, d) = \max \left\{ \begin{array}{l} \text{value}[k] + M(k-1, d - \text{size}[k]) \text{ if } k > 0, d \geq \text{size}[k] \\ M(k-1, d) \text{ if } k > 0, d \geq 0 \\ 0 \end{array} \right.$$

$4 + 10 = 14 < 15$

	size	value
0		
1	2	7
2	5	6
3	1	1
4	1	2
5	3	4
6	1	9
7	8	10
8	6	6

	d												
	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	7	7	7	7	7	7	7	7	7	7	7
2	0	0	7	7	7	7	7	13	13	13	13	13	13
3	0	1	7	8	8	8	8	13	14	14	14	14	14
4	0	2	7	9	10	10	10	13	15	16	16	16	16
5	0	2	7	9	10	11	13	14	15	16	17		
6	0	9											
7	0	9											
8	0	9											

$$M(k, d) = \max \begin{cases} \text{value}[k] + M(k-1, d - \text{size}[k]) & \text{if } k > 0, d \geq \text{size}[k] \\ M(k-1, d) & \text{if } k > 0, d \geq 0 \\ 0 & \end{cases}$$

	size	value
0		
1	2	7
2	5	6
3	1	1
4	1	2
5	3	4
6	1	9
7	8	10
8	6	6

		d												
		0	1	2	3	4	5	6	7	8	9	10	11	12
0		0	0	0	0	0	0	0	0	0	0	0	0	0
1		0	0	7	7	7	7	7	7	7	7	7	7	7
2		0	0	7	7	7	7	7	13	13	13	13	13	13
3		0	1	7	8	8	8	8	13	14	14	14	14	14
4	k	0	2	7	9	10	10	10	13	15	16	16	16	16
5		0	2	7	9	10	11	13	14	15	16			
6		0	9											
7		0	9											
8		0	9											

$$M(k, d) = \max \begin{cases} \text{value}[k] + M(k-1, d - \text{size}[k]) & \text{if } k > 0, d \geq \text{size}[k] \\ M(k-1, d) & \text{if } k > 0, d \geq 0 \\ 0 & \end{cases}$$

	size	value
0	X	X
1	2	7
2	5	6
3	1	1
4	1	2
5	3	4
6	1	9
7	8	10
8	6	6

		d												
		0	1	2	3	4	5	6	7	8	9	10	11	12
0		0	0	0	0	0	0	0	0	0	0	0	0	0
1		0	0	7	7	7	7	7	7	7	7	7	7	7
2		0	0	7	7	7	7	7	13	13	13	13	13	13
3		0	1	7	8	8	8	8	13	14	14	14	14	14
4	k	0	2	7	9	10	10	10	13	15	16	16	16	16
5		0	2	7	9	10	11	13	14	15	16	17	19	20
6		0	9	11	16	18	19	20	22	23	24	25	26	28
7		0	9	11	16	18	19	20	22	23	24	25	26	28
8		0	9	11	16	18	19	20	22	23	24	25	26	28

Back-tracing
How did I get here?

$$M(k, d) = \max \begin{cases} \text{value}[k] + M(k-1, d - \text{size}[k]) & \text{if } k > 0, d \geq \text{size}[k] \\ M(k-1, d) & \text{if } k > 0, d \geq 0 \\ 0 & \end{cases}$$

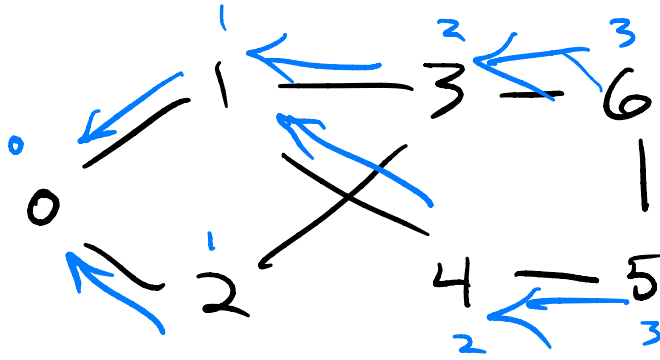
Backtracing

- Problems: maximize/minimize $F(s)$
over all $s \in S$ (search space)
- Algorithm: finds best value of $F(s)$
- Question: can you also get $m \in S$
such that $F(m)$ is optimal
(not just the optimal value, but
also the way we get it)

Backtracing

- BFS: got v 's distance from start node u , but what's the path from u to v ?
- Dijkstra: same
- Wire-cutting: best profit, but where should I make the cuts?
- Knapsack: best total value, but which objects should I select?

BFS



Queue 1 2 3 4

	0	1	2	3	4	5	6
parent?	-1	0	0	1	1	4	3
predecessor?							

	0	1	2	3	4	5	6
P	0	2	3	8	9	12	13

$$M(n) = \max_{i=1, \dots, n} (p[i] + M(n-i))$$

$$M(0) = 0$$

	0	1	2	3	4	5	6
M	0	2	4	8	10	12	16
P	0	2	3	8	9	12	13

	0	1	2	3	4	5	6
P	0	2	3	8	9	12	13

$$M(n) = \max_{i=1, \dots, n} (p[i] + M(n-i))$$

$$M(0) = 0$$

	0	1	2	3	4	5	6
M	0	2	4	8	10	12	16
P	0	2	3	8	9	12	13

	size	value
0		
1	2	7
2	5	6
3	1	1
4	1	2
5	3	4
6	1	9
7	8	10
8	6	6

	d												
	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0												
2	0												
3	0												
4	0												
5	0												
6	0												
7	0												
8	0												

$$M(k, d) = \max \begin{cases} \text{value}[k] + M(k-1, d - \text{size}[k]) & \text{if } k > 0, d \geq \text{size}[k] \\ M(k-1, d) & \text{if } k > 0, d \geq 0 \\ 0 & \end{cases}$$