

Scheme Reference Guide

<code>(define var expr)</code>	Evaluates <i>expr</i> . Binds the variable name <i>var</i> to the value in the global environment.
<code>(quote object)</code> <code>' object</code>	Returns <i>object</i> . E.g., <code>'(1 2 3)</code> returns <code>(1 2 3)</code> .
<code>(lambda (param1 ... paramn) body)</code>	Returns a procedure that takes <i>n</i> expressions as its arguments, and when invoked evaluates the <i>body</i> . For example, the value of <code>(lambda (x) (* x x))</code> is a one-argument procedure (which squares its input). To invoke this procedure, put it in parentheses with its arguments. Thus <code>((lambda (x) (* x x)) 5)</code> evaluates to 25.
<code>(lambda args body)</code> <code>(lambda (p1 ... pn . args) body)</code>	Syntax for variable-arity functions. If a single symbol is in the place of the arguments, the function takes any number of arguments (when the function is invoked, that symbol is bound to the list of arguments). A lambda expression with a dotted pair (e.g., <code>(a b . rest)</code>) requires one argument for each named symbol, aside from the last; when invoked, the list of remaining arguments are bound to the last symbol. (E.g., <code>((lambda (a . others) body) 1 2 3)</code> has <code>a</code> bound to 1 and <code>others</code> to <code>(2 3)</code> .)
<code>(if test conseq alt)</code>	Evaluates <i>test</i> . If the value is true (anything other than <code>#f</code>), then it evaluates <i>conseq</i> , whose value is the value of the <code>if</code> ; otherwise it evaluates <i>alt</i> , whose value is the value of the <code>if</code> .
<code>(cond (t1 v1) ... (tn vn))</code>	Evaluates <i>t1</i> , <i>t2</i> , ... until one is true; then evaluates the corresponding <i>vi</i> , which is the value of the entire <code>cond</code> . The last test <i>tn</i> can be replaced by <code>else</code> , in which case the value of the last body is returned whenever all previous <i>tis</i> are false.
<code>(let ((v1 e1) ... (vn en)) body)</code>	Evaluates <i>e1</i> , <i>e2</i> , ...; then binds corresponding <i>vis</i> to their values. The value of the expression is the evaluation of <i>body</i> .
<code>(let* ((v1 e1) ... (vn en)) body)</code>	Similar, but when <i>ek</i> is being evaluated, the first <i>k</i> - 1 values are bound.
<code>(letrec ((v1 e1) ... (vn en)) body)</code>	Similar, but all of <i>e1</i> , <i>e2</i> , ... are within the scope of all of the variables <i>v1</i> , <i>v2</i> , ... so <code>letrec</code> allows the definition of mutually recursive procedures.
<code>(apply f args)</code>	Invoke the function <i>f</i> with the arguments as <i>args</i> .
<code>(map f L)</code>	Apply the function <i>f</i> to each element of <i>L</i> , and collect the results in a list.
<code>(filter f L)</code>	Apply the function <i>f</i> to each element of <i>L</i> , and collect all elements <i>x</i> of <i>L</i> for which <code>(f x)</code> is true in a list.